Distributed Data Processing Environments

José Orlando Pereira

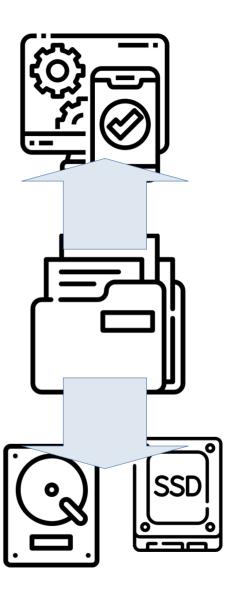
Departamento de Informática Universidade do Minho



Storage stack

What is stored in files?

How are files stored?



Motivation

- Tabular data
- Multiple data types
- Optional (null) values
- No nested or repeated values
- Large number of columns

| Id | Name | Location |
|----|--------|----------|
| 1 | aa | Braga |
| 2 | bbb | Porto |
| 3 | СС | Porto |
| 4 | dddddd | |
| 5 | eee | Lisboa |
| | | |

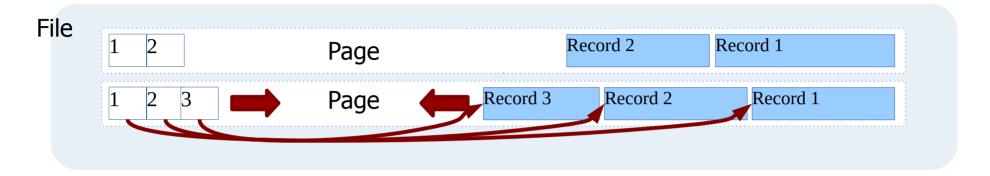
Issues

- Representation of types
 - Compactness and ambiguity
- Data that needs to be moved for:
 - Selection (range scan)
 - Projection
- Compression
- Updates

Database systems

- File format is tightly coupled to system internals
 - Optimized for direct access and caching
 - Not portable, often, even between different versions and processor architectures
- Metadata stored in schema

Row layout



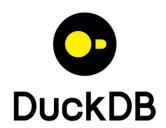
- Page based
 - Two stacks: Offset table and records
 - Fields packed in each record
- Reference for indexing:
 - (Block offset, Record index)

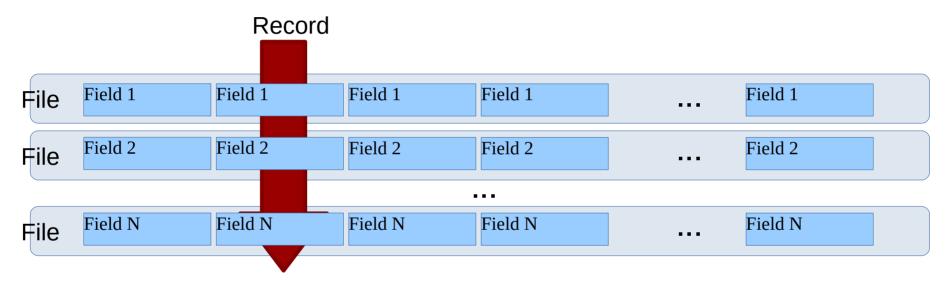


Consequences

- Efficient direct access to individual records
- I/U/D impact only one page
- Cannot avoid reading the full width of the table when scanning a column
 - e.g. select sum(a) from verylargetable;
- Good for <u>transactional</u> processing

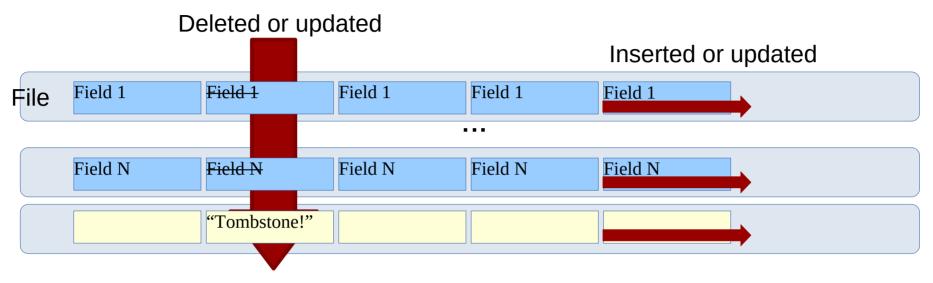
Columnar layout





 Each file holds a column (i.e., the same field for all rows)

Columnar layout



- Efficient insert (append data)
- Updates and deletes require:
 - "Tombstone" column
 - Re-writing the file

Columnar layout

- Sequences of fields of the same data type can be compressed easily and efficiently:
 - Run length encoding
 - Bitmaps and dictionaries
- Operations on compressed data improve efficiency on all levels of the memory hierarchy:



Consequences

- Column stores: Update/delete are costly operations
- Inserts are cheap by appending to files
- Efficient scan of a subset of columns from a large table
- Data already in a compact vector format
- Better for <u>analytical</u> processing

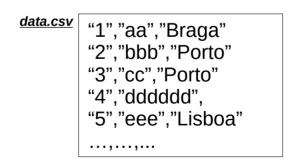
Autonomous files

- File format is decoupled from processing systems
 - Optimized for exchange and storage in a diversity of media
 - Standardized and portable
- Absent or explicitly included metadata
 - Need for external metadata management

Text (CSV)

- Simple to produce and consume
- Schema can be inferred
- Redundancy and verbose representation (numbers)
- Ambiguity in separators and missing fields
- Difficult to page, especially when compressed

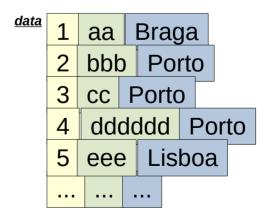
| Id | Name | Location |
|----|--------|----------|
| 1 | aa | Braga |
| 2 | bbb | Porto |
| 3 | СС | Porto |
| 4 | dddddd | |
| 5 | eee | Lisboa |
| | | |



Binary rows

- Compact and unambiguous
- Not ideal for compression, as different data types are interleaved
- All data is read for projections and selections

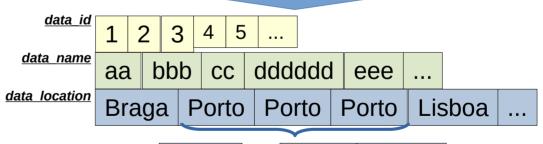
| Id | Name | Location |
|----|----------|----------|
| 1 | aa | Braga |
| 2 | bbb | Porto |
| 3 | СС | Porto |
| 4 | dddddd | |
| 5 | eee | Lisboa |
| | <u> </u> | |



Columnar

- Efficient projections
- Compressed very efficiently
 - Dictionary and/or
 - Run Length Encoding (RLE)
- Inefficient range scan

| Id | Name | Location |
|----|--------|----------|
| 1 | aa | Braga |
| 2 | bbb | Porto |
| 3 | СС | Porto |
| 4 | dddddd | |
| 5 | eee | Lisboa |
| | • | |

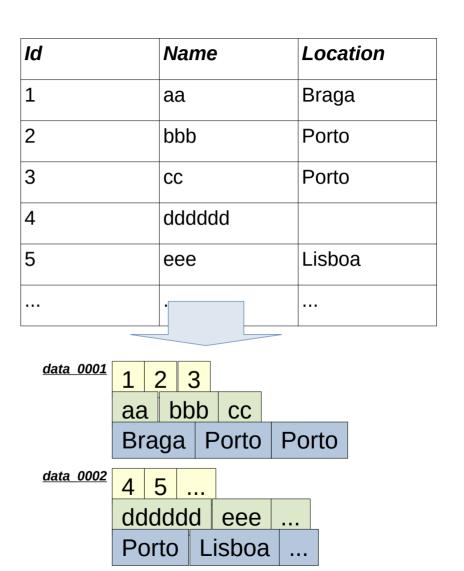


RLE: Braga 3 x Porto Lisboa

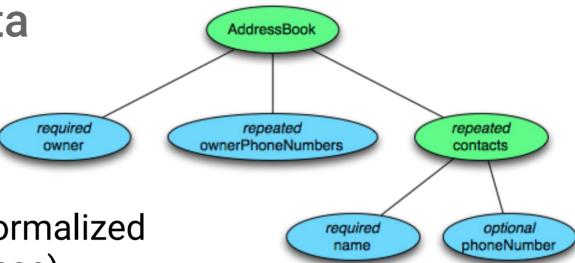
Dict.: 1 2 2 2 3 ... 1: Braga, 2: Porto, 3: Lishoa

Hybrid

- Columnar segments, that can be accessed and compressed separately
- Good trade-off:
 - I/U/D updates only one segment
 - Range scans can read only some segments
 - Projections can easily skip columns



Hierarchical data



- Data that is not normalized (in a relational sense)
 - Nested structures
 - Repeated fields
- Useful as it avoids multiple files and foreign keys

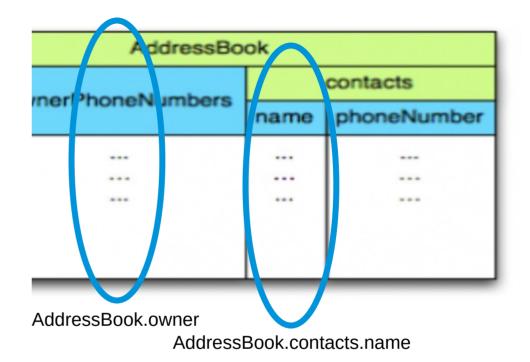
JSON

- Well-known and widely supported
- Row-based
- Not splitable

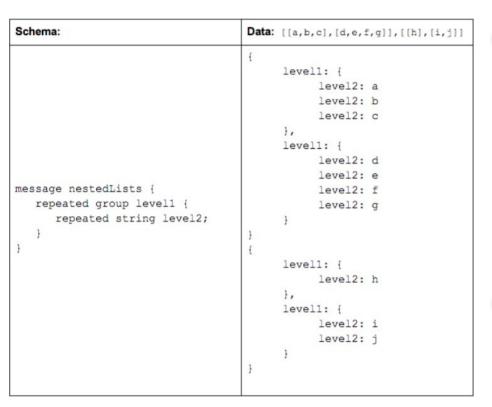
```
"AddressBook": [
  "owner": "Jason F.",
  "ownerPhoneNumbers": [
    "123456789",
    "987654321"
  "contacts": [
    { "name": "John" },
    { "name": "Joe", "number": "214365879" }
},
  "owner": "Joe G.",
  "ownerPhoneNumbers": [
    "214365879"
```

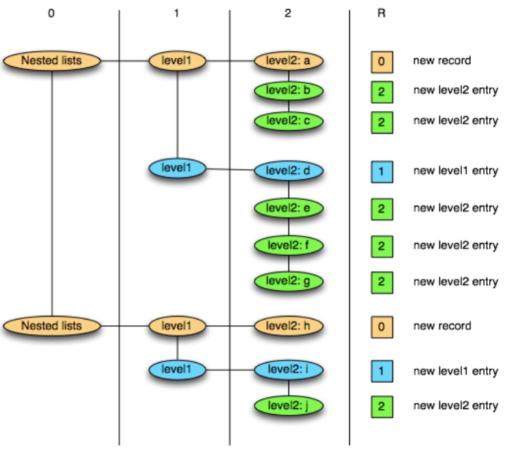
Dremel splitting

- One columnar file for each leaf attribute
- How to match records in different columns?
- Avoid additional information: record numbers, keys, ...



Dremel splitting





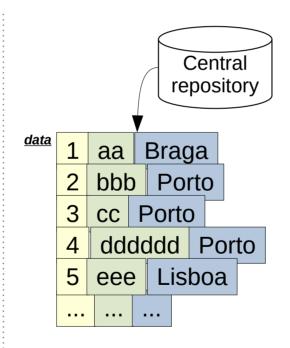
Types of metadata

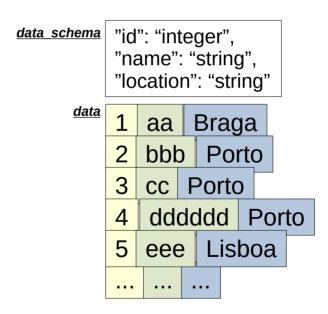
- Technical
 - Types, representation, ...
- Operational
 - Versioning, location (indexing), cardinality, ...
- Business
 - What it means, quality, ...

Schema

- Information about data items and types
- Implicit, central or embedded:





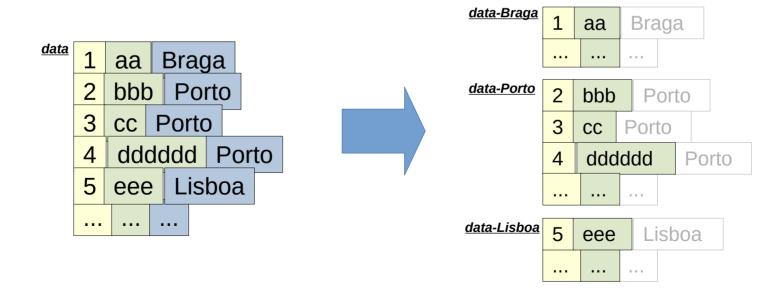


Example

```
create external table title basics (
      tconst string,
      titleType string,
      primaryTitle string,
      originalTitle string,
      isAdult boolean,
      startYear integer,
      endYear integer,
      runtimeMinutes integer,
      genres array<string>)
  row format delimited
  fields terminated by '\t'
  collection items terminated by ','
  lines terminated by '\n'
  stored as textfile
  location 'hdfs://namenode/title_basics'
  tblproperties ("skip.header.line.count"="1");
```

Partitions

- Partition files by a low cardinality column
- Encode partition key in the file name
- Used often with locations and dates
- Useful to avoid reading data

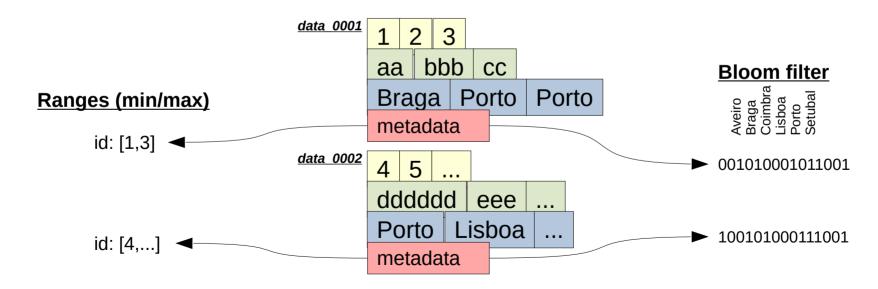


Example

```
create table title_basics_pq (
       tconst string,
       primaryTitle string,
       originalTitle string,
       isAdult boolean,
       startYear integer,
       endYear integer,
       runtimeMinutes integer,
       genres array<string>)
  partitioned by (titleType string)
  stored as parquet;
```

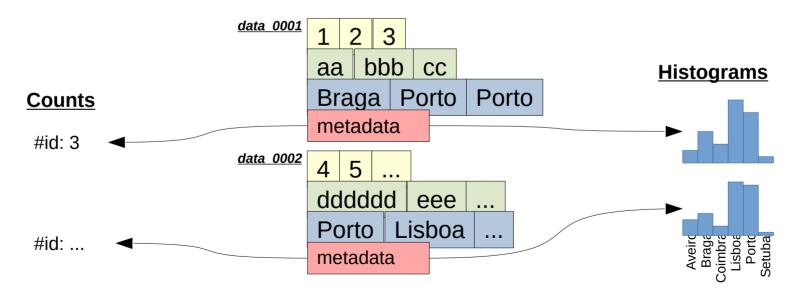
Value summaries / indexes

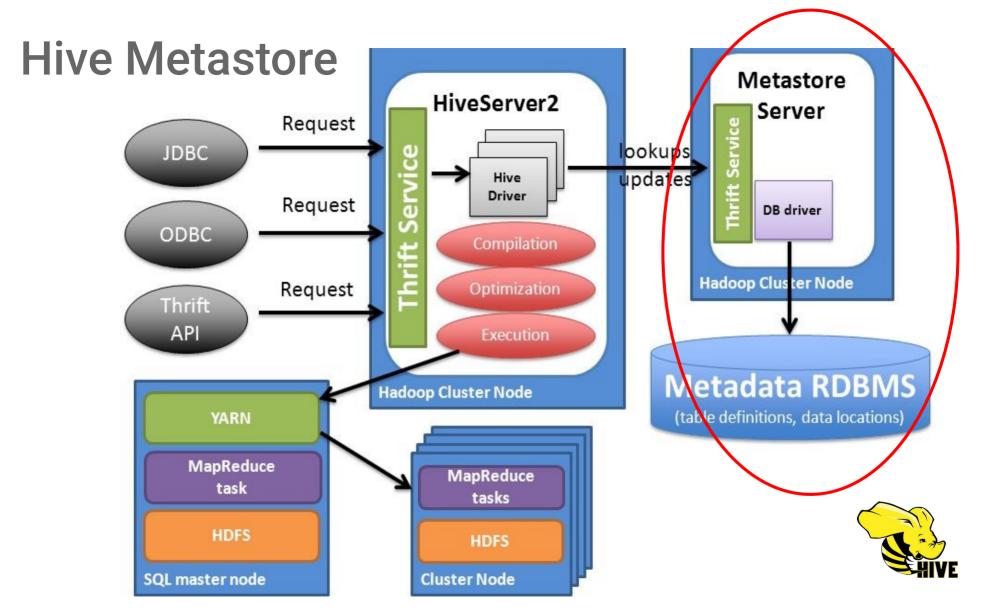
- Range [min,max] of values in each column
- Compact representation (e.g., Bloom filter) of values in each column
- Useful to avoid reading data



Cardinality summaries

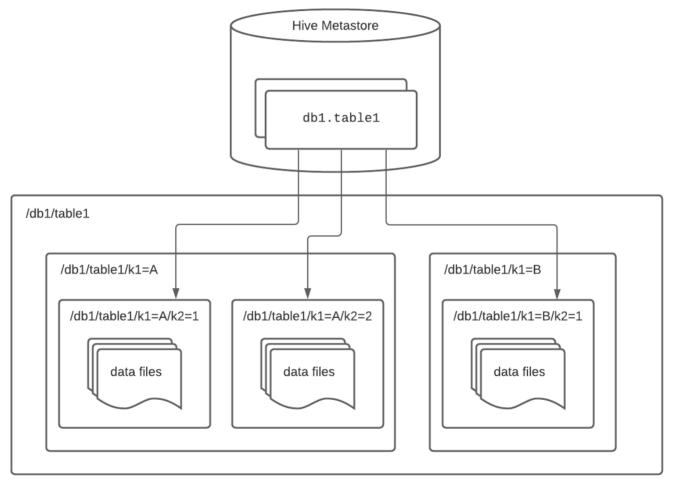
- Number of distinct values in each column
- Compact representation (e.g., histogram) of repetitions of values in intervals, for each column
- Useful to predict how much data will be processed and stored





Hive Metastore





Hive Metastore

- Central repository for technical and operational metadata
 - Individual file locations are implicitly defined by partitions
- Enables SQL query processing over data files
 - "Datalake"
- Limited support for updating files
 - Can only add partitions
 - No transactional isolation
 - Statistics can become stale

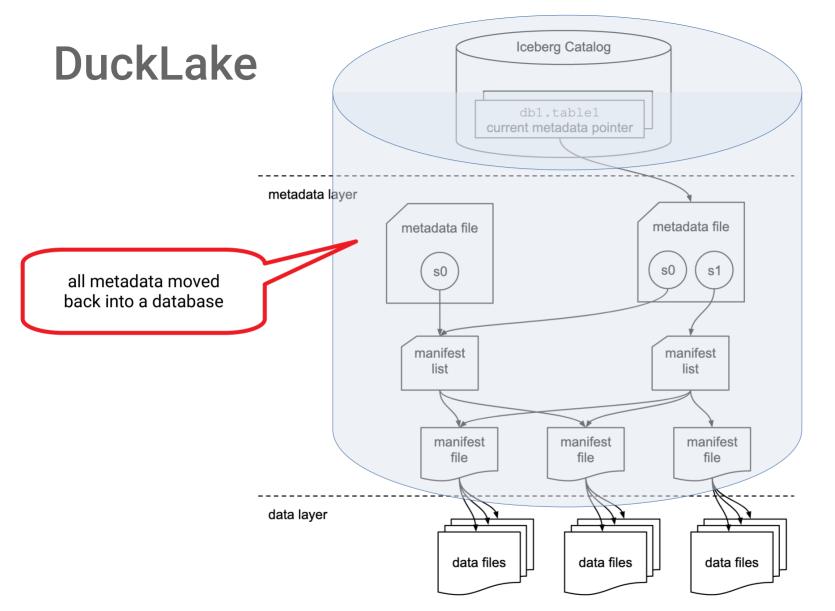
Iceberg Catalog **ICEBERG Iceberg** db1.table1 current metadata pointer metadata layer metadata file metadata file transactionally isolated updates possible by s0 s1 creating new metadata files file location is explicitly manifest manifest stored in files list list manifest manifest manifest file file file data layer data files data files data files

Iceberg

- Metadata pointer identifies the file that contains the most current information
 - Might itself be stored in a file
- A "metadata file" contains a list of snapshots, that described the file at different times
- A "manifest list" contains a list of fragments that exist at the same time
- A "manifest file" decribes an actual physical file
 - Contains statistics
- All information, except the pointer is immutable
 - No inconsistency
 - Can be stored in cloud object stores (e.g. S3)

Transactions on Iceberg

- New file fragments of new versions of existing fragments can be created by I/U/D
 - Corresponding Manifest, Manifest list, and Metadata files are created describing the new version
 - The new version is registered as current, but keeps history of previous snapshots
- Enables data modification
 - "Lakehouse"





Summary

- Minimizing data movement in analytical operations
 - Physical representation
 - Metadata
- Metadata stores key to interpreting file collections as tables
 - Datalakes
 - Lakehouses